

A security assessment study and trial of TriCore-powered automotive ECUs

CODEBLUE 2014.12.18-19

Takahiro Matsuki (FFRI)
Dennis Kengo Oka (ETAS)

Today's Talk

- ▶ Introduction
- ▶ About ECU Software
- ▶ Overview of TriCore
- ▶ Investigation and Confirmation of Attack Methods
- ▶ Demo
- ▶ Summary and Future Plans

Introduction– Motivation

- ▶ Previous research has shown that vehicle ECUs can be targeted by attackers through injection of messages on the CAN bus – what about ECU software?
- ▶ ECU microcontroller architecture is different from traditional PC architecture; therefore, traditional software attacks do not work?
- ▶ ECU microcontrollers have specific security countermeasures preventing software attacks?
- ▶ If software is vulnerable, by adjusting traditional software attacks to ECU microcontroller architecture, it is possible to execute attacks?

Knowledge Required for Security Research of ECUs

- ▶ ECU Hardware and Software Configuration
 - ECU functions and microcontrollers, bus, I/O Interface
 - What microcontrollers are used
- ▶ ECU Microcontroller Architecture
 - Program Execution Method
 - Instruction Execution Flow, Register, Memory Layout
- ▶ ECU Software Execution Environment and Development Environment
 - Library, Compiler
 - Content of the code generated by development tools
 - Reverse engineering methods of the program files

About ECU Software

Vehicle ECUs

- ▶ Basically every modern production car has a multitude of electronic control units to provide safety- as well as comfort functions.
- ▶ Average vehicle has up to 70 ECUs
- ▶ >20,000,000 lines of source code
- ▶ Electronic components are estimated to cost about 50% of the automotive production costs by 2015
- ▶ 50% infotainment, 30% powertrain and transmission, 10% chassis control, 10% body and comfort

Examples of ECUs

- ▶ Engine control unit
 - Fuel amount and mixture, air and fuel delivery timing, valve timing, ignition timing, emission control, etc...
- ▶ Transmission control unit
 - Gear change, shift lock, shift solenoids, pressure control solenoids, etc...
- ▶ Body control unit
 - Central locking, immobilizer system, power windows, climate control, etc...
- ▶ ABS/ESP control unit
 - Regulating brake pressure, traction control, cornering brake control, etc...

ECU Basic Functionality

- ▶ Typical feedback control system
- ▶ 1. Monitor input. e.g. timer, sensors, CAN
- ▶ 2. Calculate or lookup appropriate response
- ▶ 3. Generate corresponding outputs

About ECU Software

- ▶ Fully custom, proprietary software
- ▶ Unix-based proprietary software
- ▶ Standardized software. e.g. AUTOSAR

Overview of TriCore

About TriCore

- ▶ Microcontroller for Vehicle ECUs
- ▶ Manufactured and sold by Infineon
 - Spin off of Semiconductor Unit from the German manufacturer Siemens AG
- ▶ ECUs With TriCore
 - Bosch EDC17 & MED17, Siemens
- ▶ Car Manufacturers Using ECUs with TriCore
 - Audi, BMW, Citroen, Ford, Honda, Hyundai, Mercedes-Benz, Nissan, Opel, Peugeot, Porsche, Renault, Seat, Toyota, Volkswagen, Volvo

Overview of the Architecture and Features

- ▶ Command Set
 - 32 bit RISC Architecture
- ▶ Unique Register Configuration
 - Completely separated address and data registers
 - A0~A16, D0~D16
- ▶ Model Number and Specifications of the Microcontroller Used in this Research
 - TC1797 (AUDO Future)
 - TriCore Architecture 1.3.1
 - Clock 180 MHz

Research Method (1)

- ▶ Research of Open Information and Specification Documents on the Web
 - Official User's Manual
 - Most reliable information source
 - Focused on memory related charts/diagrams
 - Keyword search for security related terms
 - security, protection, password
 - TriCore Architecture Overview
 - Summarized material of the User's Manual

Research Method(2)

- ▶ Searched Research Papers
 - TriCore Emulator
 - Porting TriCore to QEMU
 - Porting Linux Kernel to TriCore
- ▶ Searched for Information/Tools for Software Developers of ECU Software
 - Development Environment TASKING VX-toolset for TriCore (Evaluation Edition)
 - Compiler, IDE with simulator
 - Evaluation Board Infineon Starter Kit TC1797
 - FlexECU development platform
- ▶ Reverse Engineering of the Binary
 - Possible to Disassemble using IDA Pro
 - File Format is ELF for Siemens TriCore



Investigation and Confirmation of Attack Methods

Possible Vulnerabilities in ECU Software

▶ Non-Memory Corruption

Vulnerabilities

- Access Control Issues
- Encryption Strength Issues
- Inappropriate Authentication
- Conflicts
- Certificate / Password Management Issues
- etc.

▶ Memory Corruption

Vulnerabilities

- Buffer Overflow
- Integer Overflow
- Use After Free
- Null Pointer Dereferences
- Format String Bugs
- etc.

Since it was difficult to obtain and analyze actual ECU Software, we hypothesized about the possibility of memory corruption vulnerabilities.

Consideration of Memory Corruption Vulnerabilities and Possible Attacks

- ▶ Buffer Overflow
 - Stack Overflow
 - Heap Overflow
- ▶ Integer Overflow
 - Hypothesized that integer overflows can cause of heap overflows
- ▶ Format String Bug
 - Possible to overwrite an arbitrary value in an arbitrary address, hypothesized that attacks are possible
- ▶ Use After Free
 - Implied attacks are possible because C++ code is executable with TriCore
- ▶ Null Pointer Dereference
 - Trap occurs by access to memory address zero, hypothesized attacks are possible

Possibility of Buffer Overflow Attacks

▶ Stack Overflows

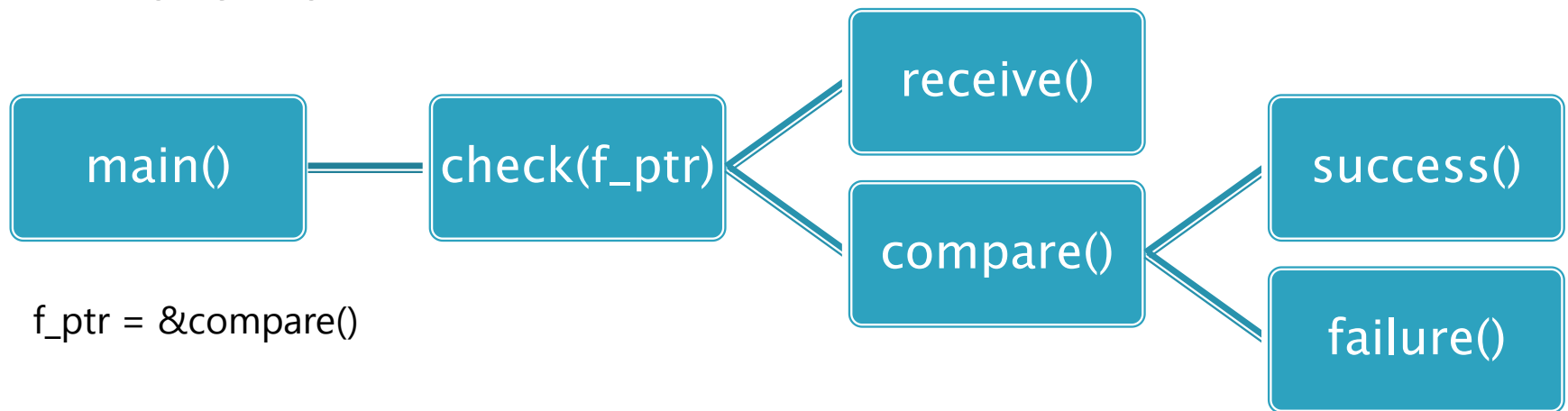
- For TriCore, unlike x86 and others, the return address is saved in the address register (A11) instead of the stack, therefore overwriting the return address using a stack overflow is not possible.

▶ Heap Overflows

- Will examine TriCore's heap management in the near future

Overwriting Function Pointers by Stack Overflows (1)

- ▶ Possibilities of a Stack Overflow Attack
 - If the buffer and function pointer exist on the stack in the following code flow, it is possible to change the execution flow of the program by a stack overflow



Overwriting Function Pointers by Stack Overflows (2)

▶ Example Code

```
failure() { ...  
}  
  
success() { ....  
}  
  
compare() { ...  
}  
  
receive() {  
    // receive input value  
    input = ...  
    char buffer[10];  
    strcpy( buffer, input );  
}
```

```
check( f_ptr ) {  
    // call receive() to receive incoming values  
    receive();  
    // call compare() using a function pointer  
    f_ptr();  
}  
  
main () {  
    function_ptr = &compare;  
    ...  
    check( function_ptr );  
}
```

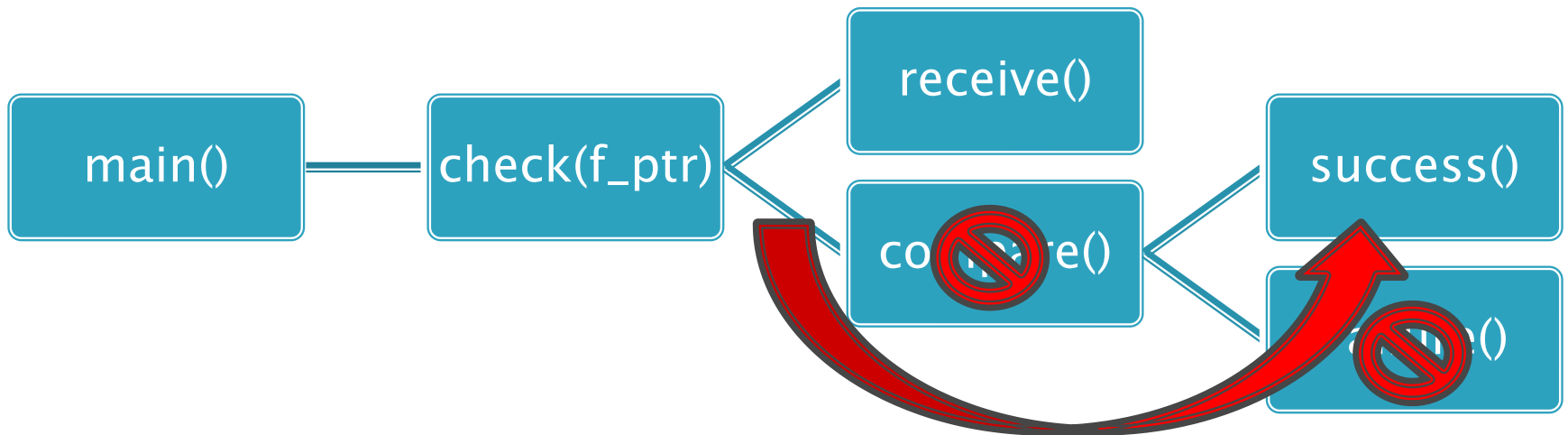
Overwriting Function Pointers by Stack Overflows (3)

► Memory Layout

Address	Variable
0x8000 010C	check()
0x8000 013C	receive()
0x8000 0170	compare()
0x8000 017C	success()
0x8000 0188	failure()
...	...
0xD000 8FC8	buffer[]
0xD000 8FE0	function_ptr



Example – Changing Flow



`check()` calls `f_ptr()` which now points to `success()` (0x8000 017C) instead of `compare()` (0x8000 0170)

Overwriting Function Pointers by Stack Overflows (4)

► Issues

- If compiler optimization is “on”, the function pointer will be stored in the address register
- Unclear whether there are similar code patterns in actual ECU software

Considering Attacks Possibilities Using TriCore's Control Mechanism

Considering Attack Possibilities Using TriCore's Control Mechanism

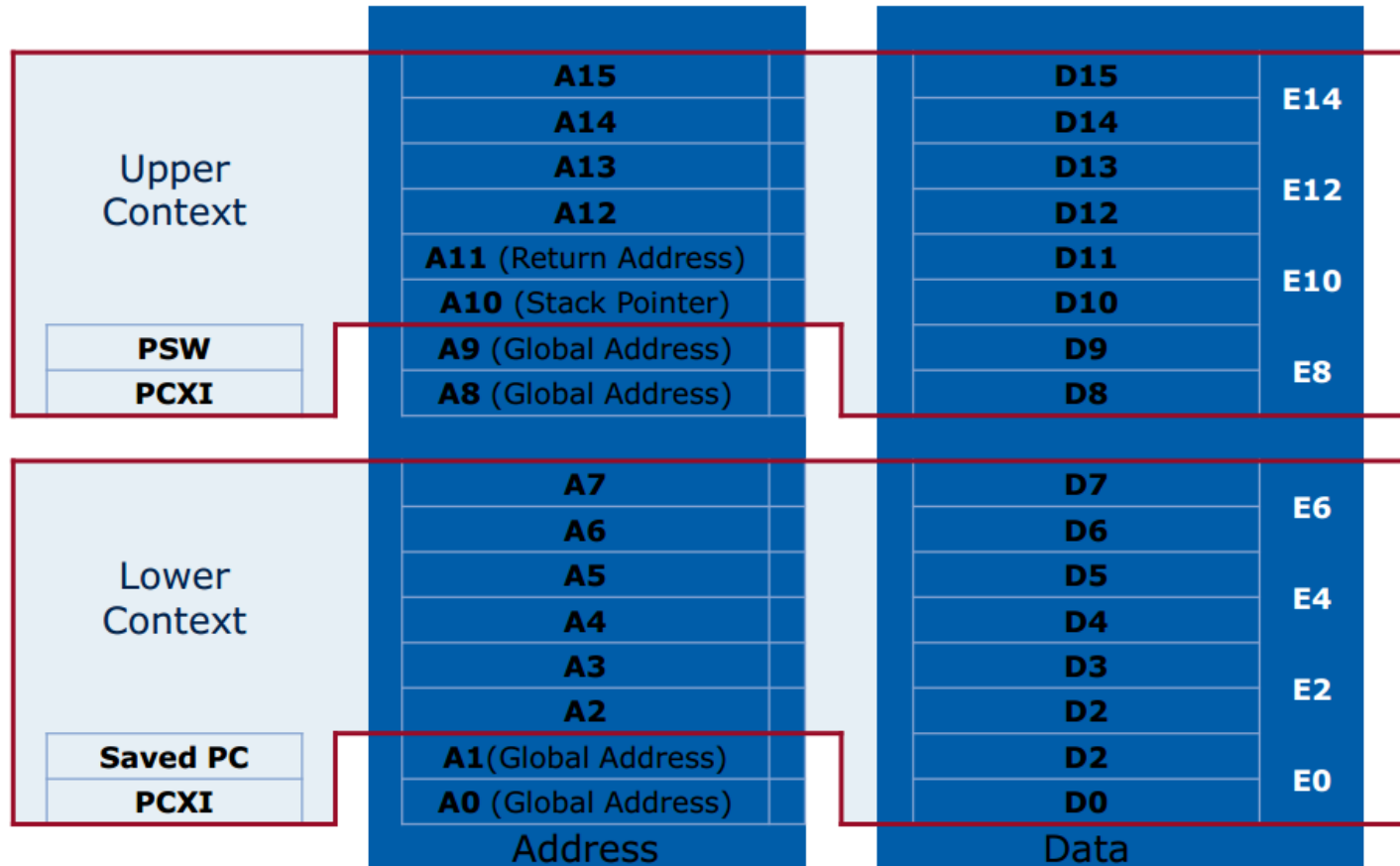
- ▶ Preconditions
 - It is possible to overwrite data by using memory corruption vulnerabilities
 - Under the above condition, considered ways to execute arbitrary code
- ▶ TriCore's Control Mechanism
 - Context Management Mechanism
 - Interrupt/Trap Mechanism

Attack Methods Using the Context Management Mechanism

Overview of Context of TriCore

- ▶ About Context
 - The register value is CSA (Context Save Area)
Saves and restores in TriCore's unique memory space
- ▶ Types of Context
 - 2 Types: Upper context and Lower context
 - Upper context
 - call command, interrupt, automatically saves when trapped
 - Lower context
 - Explicitly saved by using a dedicated command, used for passing parameters

Registers Saved in the CSA

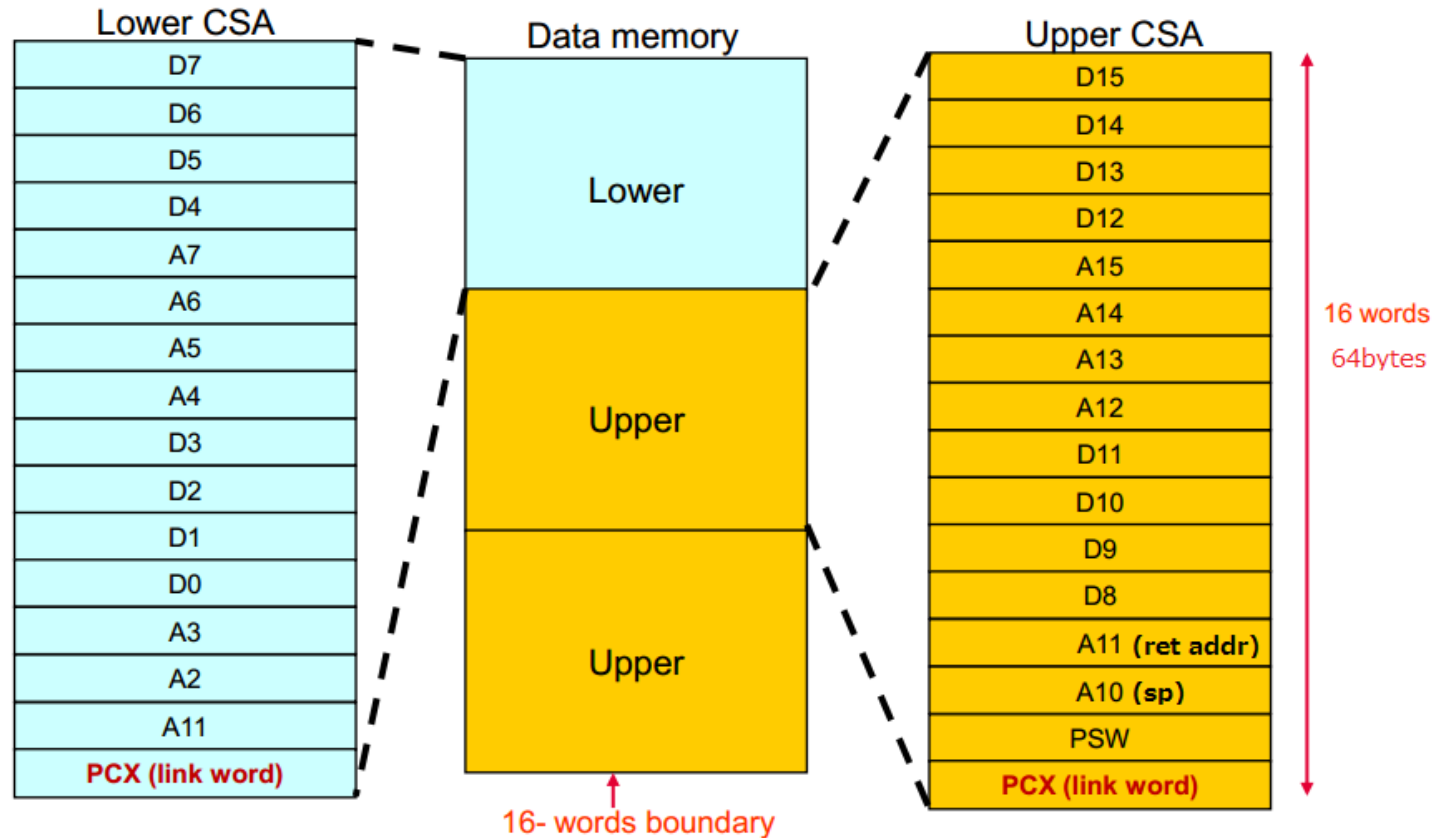


Reference: Tricore Architecture Overview

<http://www.infineon-ecosystem.org/download/schedule.php?act=detail&item=44>

CSA Configuration

- Context Save areas can hold 1 upper or 1 lower context.
- CSA are aligned on a 16-word boundary.

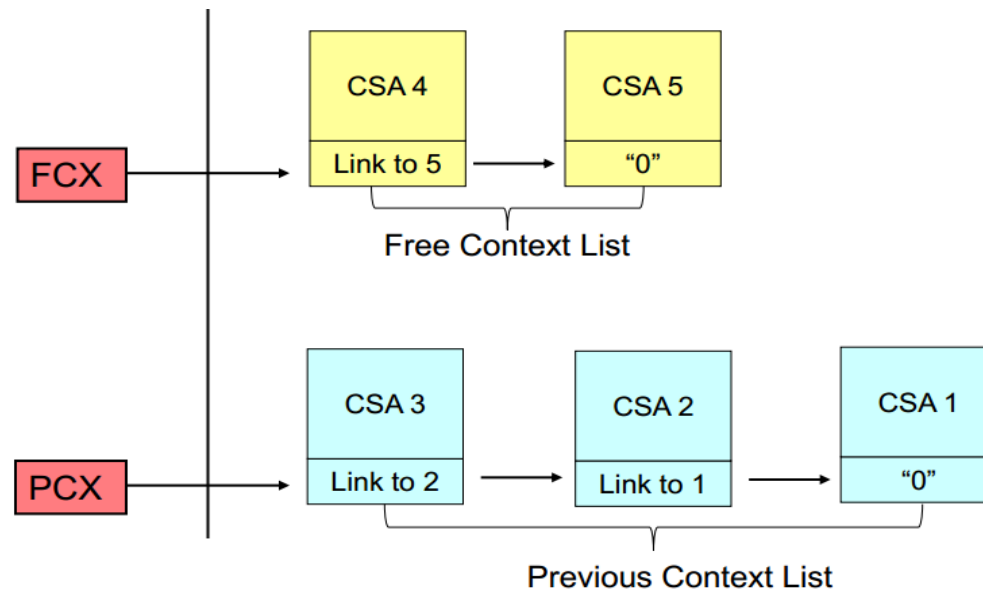


Reference: Tricore Architecture Overview

<http://www.infineon-ecosystem.org/download/schedule.php?act=detail&item=44>

CSA Management

- ▶ CSA is Managed by Link Lists
 - Used CSA List (PCX) , Unused CSA List (FCX)
 - Pointer to the first element of each list is PCX, stored in FCX register
 - However, needs to be converted because it is not a raw address



Reference: Tricore Architecture Overview

<http://www.infineon-ecosystem.org/download/schedule.php?act=detail&item=44>

Code Execution Methods Using Context

- ▶ Method 1 : CSA Overwriting
 - By overwriting any return address saved in the CSA using a memory corruption vulnerability, it is possible to run code of an arbitrary address
- ▶ Method 2 : CSA Injection
 - By overwriting a Link word of the CSA using a memory corruption vulnerability, it is possible to restore crafted Upper context (including return address) and run arbitrary code.

CSA Overwrites on a Simulator

- ▶ Code on the right is the result of execution without augments
func1
func2
func3
- ▶ Rewrite the return address (*ret) within func2 saved in the CSA to func3 address (0x80000360)
- ▶ When returned to func1, the A11 register value restores to func3 (0x80000360)
- ▶ Jump to func3 on func1 return

```
#include <stdio.h>

int func3()
{
    printf("func3\n");
    return 0;
}

int func2()
{
    unsigned int *ret;
    printf("func2\n");
    ret=0xD0004F4C;
    *ret=0x80000360;
    return 0;
}

int func1()
{
    printf("func1\n");
    func2();
    return 0;
}

int main( int argc, char** argv)
{
    func1();
    if (argc == 1)
    {
        func3();
    }
}
```

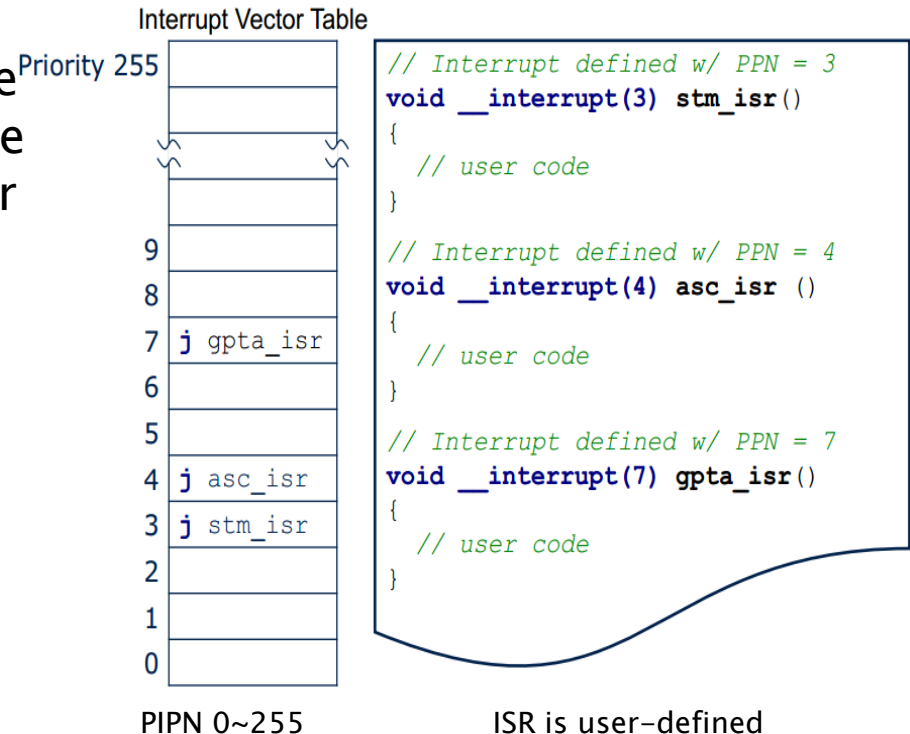
CSA Overwrites on a Simulator

- ▶ CSA overwrite using an evaluation board was possible in the same way as the simulator
 - There are memory protections to prevent CSA overwrites by default.
 - May be possible to exploit on actual ECU Software

Attack Methods Using the Interrupt and Trap Mechanisms

Overview of the Interrupt Mechanism

- ▶ When an interruption occurs, the Interrupt Vector Table (IVT) is referred, and the Interrupt Service Routine (ISR) corresponding to the Pending Interrupt Priority Number (PIPn) is executed
- ▶ IVT Start Address
 - BIV Register (Begin Interrupt Vector)
- ▶ Addresses of entry point of each ISR
 - $BIV \mid (ICR.PIPn \ll 5)$
 - ICR (Interrupt Control Register)

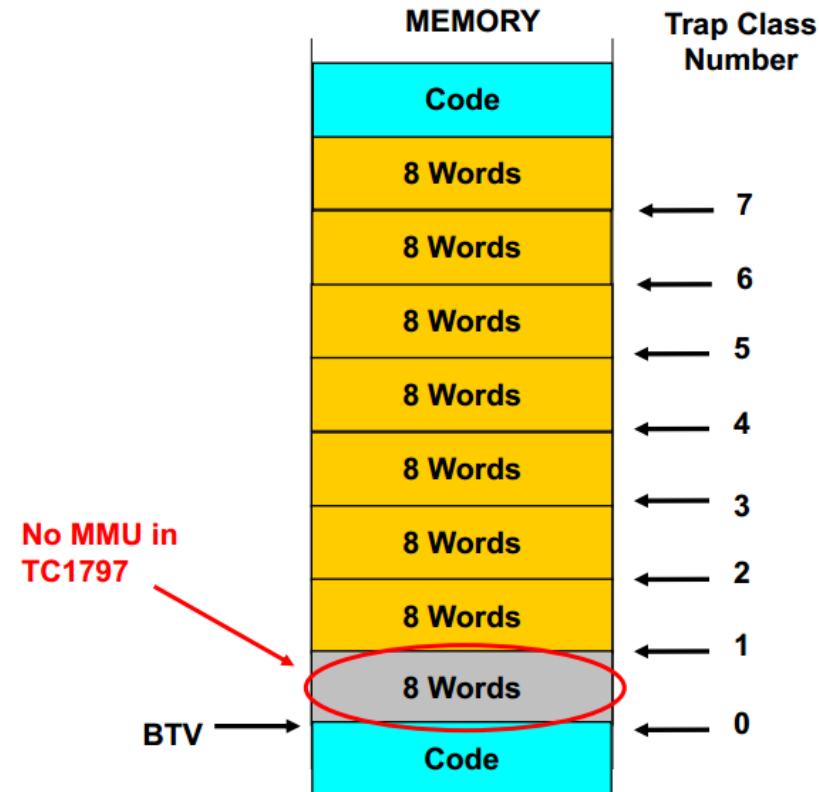


Reference: Tricore Architecture Overview

<http://www.infineon-ecosystem.org/download/schedule.php?act=detail&item=44>

Overview of the Trap Mechanism

- ▶ A mechanism used when an exception occurs. It is trapped and runs a specific process
 - Causes of Traps
 - Command exceptions, unauthorized memory access, etc...
- ▶ When a trap occurs, the Trap Vector Table is referred, and the Trap Service Routine corresponding to the Trap Class Number (TCN) is executed.



Reference: Tricore Architecture Overview

<http://www.infineon-ecosystem.org/download/schedule.php?act=detail&item=44>

Code Execution Method By Overwriting the Interrupt/Trap Vector Table

- ▶ Method 1: Overwrite the IVT
 - By overwriting the jump code to the ISR in the IVT, when a certain interrupt occurs, run arbitrary code
- ▶ Method 2: Overwrite the TVT
 - By overwriting the TSR code in the TVT , when a certain trap occurs, run arbitrary code

Overwrite of IVT/TVT on a Simulator

- ▶ BIV value 0xa00f0000
- ▶ Define a ISR as `__interrupt(3)` `hoge_isr()` , a jump code to `hoge_isr()` is allocated to 0xa00f0060 ($0xa00f00 + 32 * 3$) , making it possible to overwrite
- ▶ Overwrite possible on simulator
 - However, because whether an interrupt could be triggered intentionally is unknown, left untested
 - In the real project, the 0xA segment is mapped on the Flash memory, and may not be overwritable.
- ▶ The TVT is similarly overwritten

The screenshot displays two windows from a development environment. The top window, titled 'test.c', contains the following C code:

```
int func2()
{
    unsigned int *ret;
    printf("func2\n");
    // IVT Overwriting
    ret=0xA00F0060;
    *ret=0x01b880dd;
    return 0;
}

int func1()
{
    printf("func1\n");
    func2();
    return 0;
}

void __interrupt(3) hoge_isr()
{
    int a = 256;
    printf("int3");
}
```

The bottom window is a disassembler titled 'Disassembly'. It shows the assembly code for the address 0xa00f0060, which corresponds to the memory location where the interrupt vector is stored. The disassembly is as follows:

Address	Hex	Op	Comment
0xa00f0060	01b880dd	jla	func3 (0x80000370)
0xa00f0064	e000eed9	lea	a14, [a14] 0x380
0xa00f0068	0edc	ji	a14
0xa00f006a	0000	nop	
0xa00f006c	0000	nop	
0xa00f006e	0000	nop	

IVT/TVT Overwrite on an Evaluation Board

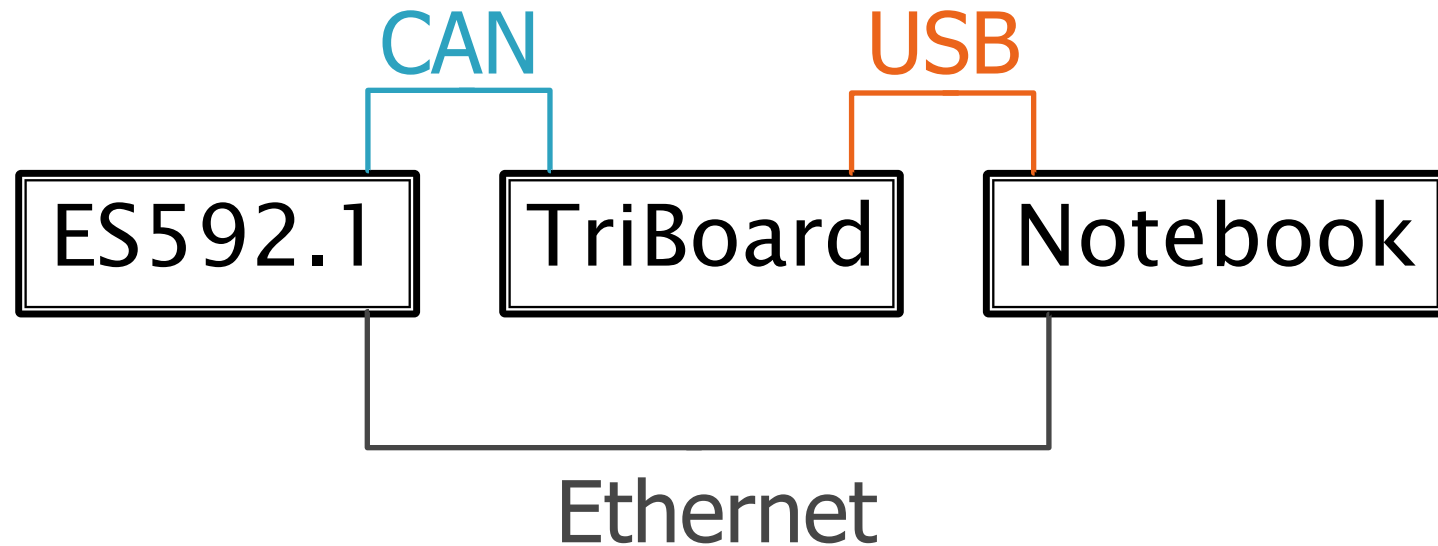
- ▶ The BIV and BTV values of the evaluation board are different from the simulator
 - BIT @ 0xF7E1FE20, BTV @ 0xF7E1FE24
- ▶ Protected
- ▶ Overwriting from the debugger possible
- ▶ Tried to overwrite the code by disabling the protection and a trap2 occurred
- ▶ Probably cannot exploit on actual ECU software

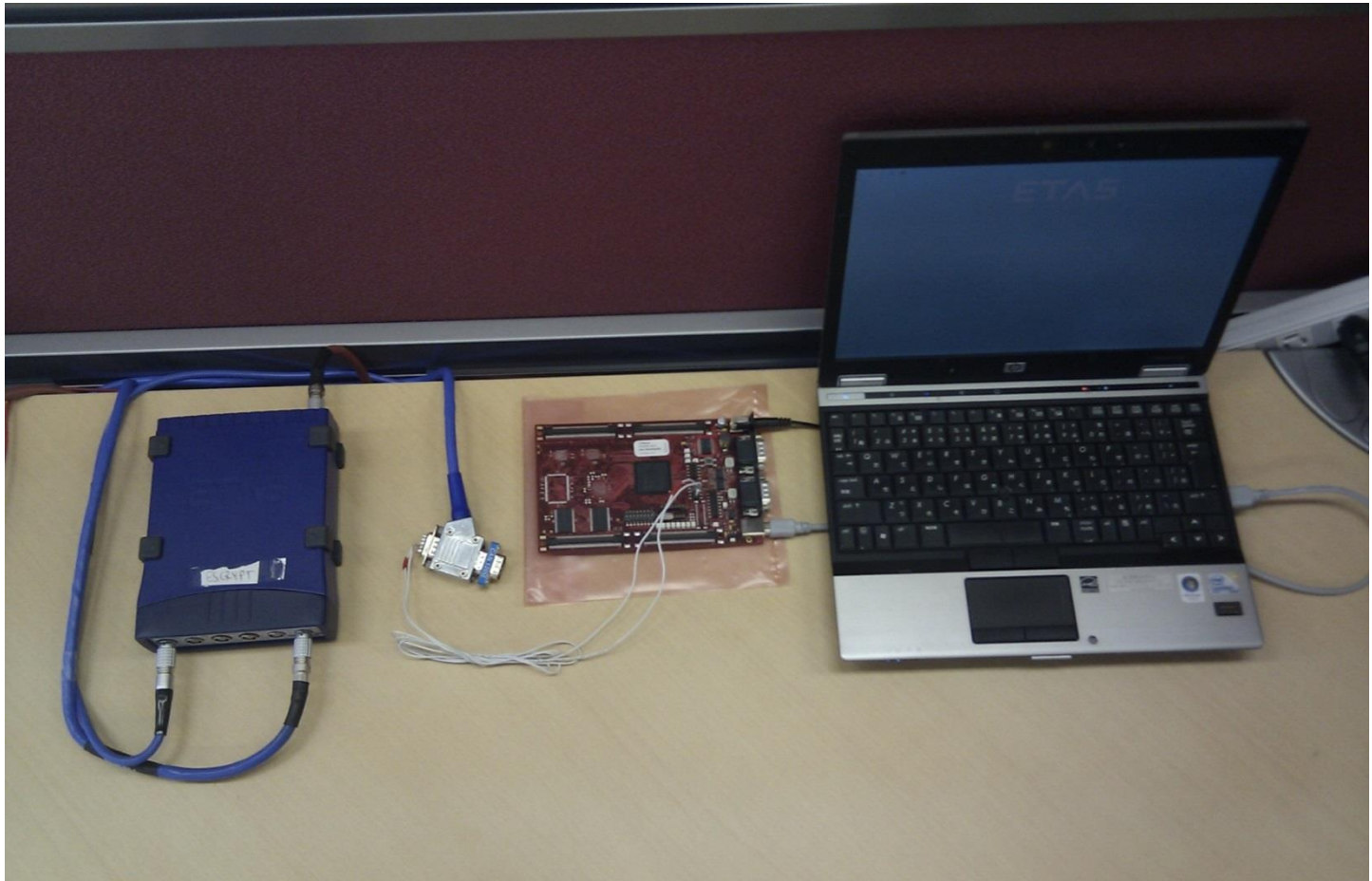
Verification on an Evaluation Board

Evaluation Environment

- ▶ HP EliteBook 2530p, Win7, Centrino2
 - HIGHTEC Free TriCore Entry Tool Chain
 - BUSMASTER
- ▶ Infineon TriBoard TC1797 V5.0
- ▶ ETAS ES592.1

Memory Monitoring Method





Demo

Summary and Future Plans

Summary

- ▶ Considered attack methods on ECU software in which memory corruption vulnerabilities exist
- ▶ If memory corruption vulnerabilities exist, it may be possible to execute arbitrary code
 - If a buffer overflow exists, it is possible to execute arbitrary code under certain conditions
 - By altering the CSA, it is possible to execute arbitrary code
 - By altering the interrupt/trap vector tables, it is impossible to execute arbitrary code
- ▶ Created vulnerable ECU software and conducted an attack demo
- ▶ This research is a result of a study of logical attack methods and a demo conducted on a vulnerable software sample. This study does NOT indicate anything about existing threats on actual ECU software.

Future Plans

▶ Additional Research

- Study other vulnerabilities and architecture specific issues

▶ Demonstrate the Threats

- Reverse engineering of ECU software and investigate if memory corruption vulnerabilities exist
- Attack actual vulnerabilities and verify if the ECUs stop or if anything abnormal occurs

▶ Consider Countermeasures

- Consider countermeasures of ECU software vulnerabilities
- Consider measures to efficiently discover vulnerabilities resulting from programming errors

Thank you